

Provably unlinkable smart card-based payments

28 November, 2023



Supported by the Luxembourg National Research Fund through grant PRIDE15/10621687/SPsquared.

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uxembourg National **Research Fund**



card number (PAN)





SMART CARD PAYMENTS (EMV)



ACTIVE ATTACKERS CAN ACTIVATE THE CARD

Passive eavesdropping

- To the card an active attacker is indistinguishable from the honest terminal

20m 1m 23 4567 8901 23 Active communication

• The cardholder, however, never enters their PIN into a random terminal that pop up on the street





- Fast
- The support of PIN
- TX:
 - Offline/Online •
 - Contact/Contactless
 - High/Low-Value

REQUIREMENTS

Security

- T authenticates C
 - T checks the legitimacy of C
 - T checks that C is not expired
- Agreement
 - If B accepts the transaction, then B, T, and C agree on the transaction

Privacy



- NO card number PAN
- NO certificate (public key, signature)
- NO expiry date



UTX PROTOCOL: PHASES





THE ESSENCE OF UTX

- Each month PaySys reveals the signed bank's public key + the validation key
- The card responds to the current (or previous) month by presenting the month certificate
- The card generates a session key with the bank and encrypts the card number PAN



VERHEUL SIGNATURES

 $\operatorname{check}(\langle M, \operatorname{vsig}(M, s) \rangle, \operatorname{vpk}(s)) = \operatorname{OK}$ $\operatorname{check}(\langle \phi(a, M), \phi(a, \operatorname{vsig}(M, s)) \rangle, \operatorname{vpk}(s)) = \operatorname{OK}$





UNLINKABILITY (DEFINITION)

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 $v \text{ user, s, si, } \chi_{MM}.\overline{out} \langle pk(s) \rangle.\overline{out} \langle vpk(\chi_{MM}) \rangle.$!vPIN, mk, c, PAN.(let crtC := $vsig(\phi(c, g), \chi_{MM})$ in vch.card(ch).C(ch, c, pk(s), crtC, PAN, mk, PIN) $|\overline{user}\langle \text{PIN}\rangle| |\langle si, \text{PAN}\rangle\langle\langle \text{PIN}, mk, \phi(c, \mathfrak{g})\rangle\rangle\rangle|$ vb_t.!vkbt. $vch.bank(ch).B(ch, si, kbt, b_t)$ let crt := $\langle \langle \mathsf{MM}, \phi(b_t, \mathfrak{g}) \rangle$, sig $(\langle \mathsf{MM}, \phi(b_t, \mathfrak{g}) \rangle, s) \rangle$ in $vch.\overline{term}\langle ch \rangle.T(user, ch, vpk(\chi_{MM}), crt, kbt)$

A card can participate in many sessions.



v user, s, si, χ_{MM} . \overline{out} $\langle pk(s) \rangle$. \overline{out} $\langle vpk(\chi_{MM}) \rangle$. !*v*PIN, *mk*, *c*, PAN.(let crtC := $vsig(\phi(c, g), \chi_{MM})$ in vch.card(ch).C(ch, c, pk(s), crtC, PAN, mk, PIN) $|\overline{user}\langle \text{PIN}\rangle| |\langle si, \text{PAN}\rangle\langle\langle \text{PIN}, mk, \phi(c, \mathfrak{g})\rangle\rangle |$ vb_t.!vkbt. $vch.bank(ch).B(ch, si, kbt, b_t)$ let crt := $\langle \langle \mathsf{MM}, \phi(b_t, \mathfrak{g}) \rangle$, sig $(\langle \mathsf{MM}, \phi(b_t, \mathfrak{g}) \rangle, s) \rangle$ in $vch.\overline{term}\langle ch \rangle.T(user, ch, vpk(\chi_{MM}), crt, kbt)$

A card can participate in at most one session.



UNLINKABILITY (PROOF CERTIFICATE)

 $(\vec{K}, F, A, \Gamma, B, \Lambda)_{\text{impl}}(X, Y, Z) \triangleq v\vec{\epsilon}, \text{PIN}_{1...H}, mk_{1...H}, c_{1...H}, \text{PAN}_{1...H}, \dot{ch}_{1...D},$ $a_{1...E}, b_t, \ddot{c}h_{1...F+G}, \ddot{c}h_{1...F+M}$ $t_{1...L}, \mathsf{TX}_{1...L}.(\theta \mid$ $C_1^1 \mid U_1^1 \mid DB_1^1 \mid$ $C_{i_1}^1 \mid U_{i_1}^1 \mid DB_{i_1}^1 \mid$ $C_{D_1+K_1}^1 \mid U_{D_1+K_1}^1 \mid DB_{D_1+K_1}^1 \mid$ $!(vch.card\langle ch \rangle.$ $C(ch, c_j, pk(s), vsig(\phi(c, g), \chi_{MM}), PAN_j, mk_j, PIN_j) |$ \overline{user} (PIN₁) | $DB(si, PAN_1, mk_1, PIN_1))$ | $C_{D_{h-1}+K_{h-1}+1}^{h} \mid U_{D_{h-1}+K_{h-1}+1}^{h} \mid DB_{D_{h-1}+K_{h-1}+1}^{h} \mid$ $C_{i_h}^h \mid U_{i_h}^h \mid DB_{i_h}^h \mid$ $C_{D_{h-1}+K_{h-1}+D_{h}+K_{h}}^{h} \mid U_{D_{h-1}+K_{h-1}+D_{h}+K_{h}}^{h} \mid$ $DB_{D_{h-1}+K_{h-1}+D_h+K_h}^h \mid$ $!(vch.card\langle ch \rangle.$ $C(ch, c_h, pk(s), vsig(\phi(c, g), \chi_{MM}), PAN_h, mk_h, PIN_h) |$ $\overline{user} \langle \text{PIN}_h \rangle | DB(si, \text{PAN}_h, mk_h, \text{PIN}_h)) |$ $C_{D_{H-1}+K_{H-1}+1}^{H} \mid U_{D_{H-1}+K_{H-1}+1}^{H} \mid DB_{D_{H-1}+K_{H-1}+1}^{H} \mid$ $C_{i_{H}}^{H} \mid U_{i_{H}}^{H} \mid DB_{i_{H}}^{H} \mid$ $C_{D_{H-1}+K_{H-1}+D_{H}+K_{H}}^{H} \mid U_{D_{H-1}+K_{H-1}+D_{H}+K_{H}}^{H} \mid$ $DB_{D_{H-1}+K_{H-1}+D_{H}+K_{H}}^{H} \mid$ $!(vch.card\langle ch \rangle.$ $C(ch, c_H, \mathsf{pk}(s), \mathsf{vsig}(\phi(c, \mathfrak{g}), \chi_{\mathsf{MM}}), \mathsf{PAN}_H, mk_H, \mathsf{PIN}_H) \mid$ \overline{user} (PIN_H) | $DB(si, PAN_H, mk_H, PIN_H))$ | !PC_{impl} | $B_1^\theta \mid T_1^\theta \mid$... $B_{j}^{\theta} \mid T_{j}^{\theta}$ $B_{F+G+M}^{\theta} \mid T_{F+G+M}^{\theta} \mid !PBT)$

 $(K, F, A, \Gamma, B)_{\text{spec}}(X, Y, Z) \triangleq$ $v \vec{\epsilon}$, PIN_{1...D+K}, $mk_{1...D+K}$, $c_{1...D+K}$, PAN_{1...D+K}, $\dot{ch}_{1...D}, a_{1...E}, b_t, \ddot{ch}_{1...F+G},$ $\ddot{ch}_{1...F+M}, t_{1...L}, \mathsf{TX}_{1...L}.(\sigma \mid$ $C_1 \mid \ldots \mid 0 \mid !\overline{user} \langle \text{PIN}_1 \rangle \mid$... $| 0 | ! \overline{\langle si, PAN_1 \rangle} \langle \langle PIN_1, mk_1, \phi(c_1, \mathfrak{g}) \rangle \rangle |$

 $C_i \mid \ldots \mid 0 \mid !\overline{user} \langle \text{PIN}_i \rangle \mid$ R

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C_{D+K} \mid \ldots \mid 0 \mid !\overline{user} \langle \text{PIN}_{D+K} \rangle \mid
     ... |0| |\langle si, PAN_{D+K} \rangle \langle \langle PIN_{D+K}, mk_{D+K}, \phi(c_{D+K}, \mathfrak{g}) \rangle \rangle |
!PC<sub>spec</sub> |
B_1^{\sigma} \mid T_1^{\sigma} \mid
...|
B_j^{\sigma} \mid T_j^{\sigma} \mid
...|
B_{F+G+M}^{\sigma} \mid T_{F+G+M}^{\sigma} \mid !PBT)
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... $| 0 | ! \overline{\langle si, PAN_i \rangle} \langle \langle PIN_i, mk_i, \phi(c_i, \mathfrak{g}) \rangle \rangle |$

 \mathfrak{R} is a quasi-open bisimulation:





CONCLUSION

- Privacy-preserving smart card payments are feasible
 - UTX is unlinkable in the presence of active attackers
 - UTX respects the essential security guarantees card payments provide
 - UTX requires only a software update to the current payment infrastructure
 - UTX can coexist with traditional card payments
- It is feasible to prove bisimilarity-based properties of complex protocols



